### Tiny functions for lots of things

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 $\blacktriangleright$  A little "functional-ish" programming goes a long way.

 $\blacktriangleright$  It's worth refactoring megamodules (codecs, TCP, compilers, machine learning) using ideas from functional programming.

 $\blacktriangleright$  Just the ability to name, save, and restore program states is powerful in its own right.

### Breaking megamodules into functions

Lepton: JPEG recompression in a distributed filesystem

ExCamera: Fast interactive video encoding

Salsify: Videoconferencing with co-designed codec and transport protocol

gg: IR for "laptop to lambda" jobs with 8,000-way parallelism

### Breaking megamodules into functions

Lepton: JPEG recompression in a distributed filesystem

 $\blacktriangleright$  "functional" JPEG codec for boundary-oblivious sharding

### ExCamera: Fast interactive video encoding

 $\blacktriangleright$  "functional" video codec for fine-grained **parallelism** 

Salsify: Videoconferencing with co-designed codec and transport protocol

 $\blacktriangleright$  "functional" codec to explore an execution path without committing

gg: IR for "laptop to lambda" jobs with 8,000-way parallelism

 $\blacktriangleright$  "functional" representation of **practical parallel pipelines** 

### System 1: Lepton (distributed JPEG recompression)

Daniel Reiter Horn, Ken Elkabany, Chris Lesniewski-Lass, and KW, The Design, Implementation, and Deployment of a System to Transparently Compress Hundreds of Petabytes of Image Files for a File-Storage Service, in NSDI 2017 (Community Award winner).

### **Storage Overview at Dropbox**

 $\cdot$   $\frac{3}{4}$  Media



### • Roughly an Exabyte in storage

• Can we save backend space?

### **JPEG File**

- Header
- 8x8 blocks of pixels
	- DCT transformed into 64 coefs

o Lossless

- Each divided by large quantizer o Lossy
- Serialized using Huffman code
	- o Lossless



*Image credit: wikimedia*



Idea: save storage with transparent recompression

 $\triangleright$  **Requirement:** byte-for-byte reconstruction of original file

▶ Approach: improve bottom "lossless" layer only

▶ Replace DC-predicted Huffman code with an arithmetic code

 $\triangleright$  Use a probability model to predict "1" vs. "0"

### Prior work



### Challenge: distributed filesystem with arbitrary chunk boundaries



### Challenge: distributed filesystem with arbitrary chunk boundaries



### Challenge: distributed filesystem with arbitrary chunk boundaries



### Requirements for distributed compression

### $\triangleright$  Store and decode file in independent chunks

 $\blacktriangleright$  Can start at any byte offset

### $\triangleright$  Achieve  $> 100$  Mbps decoding speed per chunk

### $\blacktriangleright$  Don't lose data

- $\blacktriangleright$  Immune to adversarial/pathological input files
- $\triangleright$  Every time program changed, qualify on a billion images
- $\triangleright$  Three compilers (with and without sanitizers) must match on all billion images



 $\triangleright$  Baseline JPEG is encoded as a *stream* of Huffman codewords with opaque state (DC prediction).

▶ encode(HuffmanTable, vector<Coefficient>)  $\rightarrow$  vectorshit>

- $\blacktriangleright$  How to encode chunk of original file, starting in midstream?
	- $\blacktriangleright$  Midstream  $\equiv$  in the middle of a Huffman codeword
	- $\triangleright$  Midstream = unknown DC (average) value

### When the client retrieves a chunk of a JPEG file, how does the fileserver re-encode that chunk from Lepton back to JPEG?

### Making the state of the JPEG encoder explicit

 $\blacktriangleright$ 

### $\triangleright$  Formulate JPEG encoder in explicit state-passing style

 $\blacktriangleright$  Implement DC-predicted Huffman encoder that can resume from any byte boundary

encode(HuffmanTable, vector<br/>bit>, int dc, vector<Coefficient>)  $\rightarrow$  vectorshit>

### **Results**



### **Results**



### **Deployment**

- Lepton has encoded 150 billion files
	- 203 PiB of JPEG files
	- Saving 46 PiB
	- So far…
		- o Backfilling at > 6000 images per second



### **Power Usage at 6,000 Encodes**





 $\triangleright$  A little bit of functional programming can go a long way.

▶ Functional JPEG codec lets Lepton **distribute** decoding with arbitrary chunk boundaries and **parallelize** within each chunk.

### System 2: ExCamera (fine-grained parallel video processing)

Sadjad Fouladi, Riad S. Wahby, Brennan Shacklett, Karthikeyan Vasuki Balasubramaniam, William Zeng, Rahul Bhalerao, Anirudh Sivaraman, George Porter, and KW, Encoding, Fast and Slow: Low-Latency Video Processing Using Thousands of Tiny Threads, in NSDI 2017.

<https://ex.camera>

# What we currently have

# eg Google Docs

- People can make changes to a word-processing document
- The changes are instantly visible for the others



- People can interactively edit and transform a video
- The changes are instantly visible for the others

# What we would like to have

# *for Video*?

**"Apply this awesome filter to my video."**





# **"Remake Star Wars Episode I without Jar Jar."**

# Can we achieve interactive collaborative video editing by using massive parallelism? *The Question*

# Currently, running such pipelines on videos takes hours and hours, even for a short video. *The Problem*

# The challenges

## • Low-latency video processing would need **thousands of threads**, **running in**

- **parallel**, with **instant startup.**
- **efficiency.**

### • However, **the finer-grained the parallelism, the worse the compression**





# Enter *ExCamera*

### • Framework to run **5,000-way parallel jobs** with IPC on a commercial

- We made two contributions:
	- "cloud function" service.
	-
- We call the whole system **ExCamera**.

# • Purely functional video codec for **massive fine-grained parallelism**.



# Cloud function services have (as yet) unrealized power

- AWS Lambda, Google Cloud Functions
- Intended for event handlers and Web microservices, but...
- · Features:
	- $\triangledown$  Thousands of threads
	- V Arbitrary Linux executables
	- ✔ Sub-second startup
	-

## $\checkmark$  Sub-second billing  $\checkmark$  3,600 threads for one second  $\to$  9 $\&$



### • With the existing encoders, the finer-grained the parallelism, the worse the

# Now we have the threads, but...

compression efficiency.





# Video Codec

### • A piece of software or hardware that compresses and decompresses digital

1011000101101010001 0001111111011001110 0110011101110011001 0010000...001001101 0010011011011011010 1111101001100101000 0010011011011011010 Encoder <sup>@@10000</sup>...<sup>001001101</sup> Decoder





video.



19

# How video compression works

- Exploit the temporal redundancy in adjacent images.
- Store the first image on its entirety: a **key frame**.
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• For other images, only store a "diff" with the previous images: an **interframe**.



### In a 4K video @15Mbps, a key frame is **~1 MB**, but an interframe is **~25 KB**.



# Existing video codecs only expose a simple interface

# $encode([l], l], \ldots, l] ) \rightarrow keyframe + interframe[2:n]$

21

# *compressed video*

# $decode(keyframe + interframe[2:n]) \rightarrow [\mathbf{M},\mathbf{M},\dots,\mathbf{M}]$



# Traditional parallel video encoding is limited

# $encode(i[1:200]) \rightarrow keyframe_1 + interframe[2:200]$

parallel  $\downarrow$  ————————————————————

- [thread 01] **encode**(i[1:10]) [thread 02] **encode**(i[11:20] [thread 03] **encode**(i[21:30]
- ⠇ [thread 20] **encode**(i[191:20

finer-grained parallelism  $\Rightarrow$  more key frames  $\Rightarrow$  worse compression efficiency



serial ↓

$$
\begin{array}{ccc}\n & \rightarrow & k f_1 + \text{ if } [2:10] \\
& \rightarrow & k f_1 \frac{1 \text{ MB}}{1} \text{ if } [12:20] \\
& \rightarrow & k f_2 \frac{1 \text{ MB}}{1} \text{ if } [22:30] \\
& \rightarrow & k f_1 \frac{1 \text{ MB}}{1} \text{ if } [192:200]\n\end{array}
$$
### We need a way to start encoding mid-stream

- Start encoding mid-stream needs access to intermediate computations.
- Traditional video codecs *do not* expose this information.
- We formulated this internal information and we made it explicit: the **"state"**.



### The decoder is an automaton





### The state is consisted of reference images and probability models





### What we built: a video codec in explicit state-passing style

- VP8 decoder with no inner state:
	- **decode**(state, frame) → (state′, image)
- VP8 encoder: resume from specified state **encode**(state, image) → interframe
- Adapt a frame to a different source state
	- **rebase**(state, image, interframe) → interframe′
- 



### Putting it all together: ExCamera

- Divide the video into tiny chunks:
	- [Parallel] **encode** tiny independent chunks.
	- [Serial] **rebase** the chunks together and remove extra keyframes.





### 1. [Parallel] Download a tiny chunk of raw video





### 2. [Parallel] vpxenc → keyframe, interframe[2:n]







### 3. [Parallel] decode → state → next thread





### 4. [Parallel] *last thread's state*  $\rightarrow$  encode





### 5. [Serial] *last thread's state* → rebase → state → next thread





### 5. [Serial] *last thread's state* → rebase → state → next thread



### 6. [Parallel] Upload finished video







### Wide range of different configurations



# ExCamera[**n**, **x**]

number of frames in each chunk



### Wide range of different configurations



# ExCamera[**n**, **x**] number of chunks "rebased" together

### How well does it compress?





### How well does it compress?





### 5 10 20 30 40 50 70 average bitrate (Mbit/s)

vpx (1 thread)

vpx (multithreaded)

### ExCamera[6, 1]

### How well does it compress?





### 5 10 20 30 40 50 70 average bitrate (Mbit/s)

Cameral 16

### ExCamera[6, 1]

vpx (1 thread) ±3%

# ExCamera[6, 16] **2.6 mins**

## **14.8**-minute **4K** Video *@20dB*

# vpxenc Single-Threaded **453 mins**

# vpxenc Multi-Threaded **149 mins**

# YouTube (H.264) **37 mins**

### ExCamera concluding thoughts

 $\triangleright$  Functional video codec lets ExCamera **parallelize** at fine granularity.

- $\triangleright$  Many interactive jobs call for similar approach:
	- $\blacktriangleright$  Image and video filters
	- $\blacktriangleright$  3D artists
	- $\triangleright$  Compilation and software testing
	- $\blacktriangleright$  Interactive machine learning
	- $\blacktriangleright$  Database queries
	- $\blacktriangleright$  Data visualization
	- $\blacktriangleright$  Genomics
	- $\blacktriangleright$  Search

 $\triangleright$  Distributed systems will need to treat application state as a first-class object.

Every program soon: do in 1 hour do in 1 second for  $9¢$ 

Sadjad Fouladi, John Emmons, Emre Orbay, Catherine Wu, Riad S. Wahby, and KW, Salsify: low-latency network video through tighter integration between a video codec and a transport protocol, in NSDI 2018.

<https://snr.stanford.edu/salsify>











### Current systems do not react fast enough to **network variations**, end up congesting the network, causing **stalls and glitches**.







### **Today's systems combine two** *(loosely-coupled)* **components**

10

### **Two distinct modules, two separate control loops**

11





### **Transport tells us how big the next frame should be, but...**

It's challenging for any codec to choose the appropriate quality settings upfront to meet a **target size**—they tend to over-/undershoot the target.



### **How to get an accurate frame out of an inaccurate codec**

### • **Trial and error:** Encode with different quality settings, pick the one that fits.

- - *Not possible with existing codecs.*





### **After encoding a frame, the encoder goes through a state transition that is impossible to undo**



### **There's no way to undo an encoded frame in current codecs**

22



The state is internal to the encoder—no way to save/restore the state.



### **Functional video codec to the rescue**



### **encode**(*state*, ) → *state*′ , frame



### Salsify's functional video codec exposes the state that can be saved/restored.

### **Order two, pick the one that fits!**

- without committing to them.
- For each frame, codec presents the transport with *three* options:
	- A slightly-higher-quality version,
	- A slightly-lower-quality version,
	- Discarding the frame.



### • Salsify's functional video codec can explore different execution paths







### Salsify's architecture: **Unified control loop**





### **transport protocol & video codec**

### Codec → Transport **"Here's two versions of the current frame."**







### $Transport \rightarrow Codec$ **"I picked option 2. Base the next frame on its exiting state."**









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### Codec → Transport **"Here's two versions of the latest frame."**











### Transport → Codec **"I picked option 1. Base the next frame on its exiting state."**











### Codec → Transport **"Here's two versions of the latest frame."**








# $Transport \rightarrow Codec$ **"I cannot send any frames right now. Sorry, but discard them."**







31





# Codec → Transport **"Fine. Here's two versions of the latest frame."**





# Transport → Codec **"I picked option 1. Base the next frame on its exiting state."**







## **Goals for the measurement testbed**

- A system with **reproducible input video** and **reproducible network traces** that runs **unmodified** version of the system-under-test.
- Target QoE metrics: per-frame **quality** and **delay**.



### **barcoded video**

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نبن

### **video in/out (HDMI)**

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*<u>REGISTERED</u>* 







### **receiver HDMI output**

 $\begin{array}{c|c|c|c|c} \hline \textbf{G} & \textbf{G}$ 





### **emulated network**

### **Sent Image**  Timestamp: *T+0.000s*





**Received Image**  Timestamp: *T+0.765s* Quality: *9.76* dB SSIM





41











### Evaluation results: **AT&T LTE Trace**











# Evaluation results: **T-Mobile UMTS Trace**











Improvements to *video codecs* may have reached the point of diminishing returns, but changes to the architecture of *video systems* can still yield significant benefits.

- ▶ Kalev Alpernas, Cormac Flanagan, Sadjad Fouladi, Leonid Ryzhyk, Mooly Sagiv, Thomas Schmitz, and KW, Secure serverless computing using dynamic information flow control, Proc. ACM Program. Lang. 2, OOPSLA, Article 118 (November 2018).
- ▶ Sadjad Fouladi, Francisco Romero, Dan Iter, Qian Li, Shuvo Chatterjee, Christos Kozyrakis, Matei Zaharia, and KW, From Laptop to Lambda: Outsourcing Everyday Jobs to Thousands of Transient Functional Containers, in USENIX ATC 2019.

### Cloud functions as a new computing substrate

Rent 8,000 nodes in seconds (but some are flaky)

 $\triangleright$  Nodes can communicate directly at 600 Mbps (but some paths are flaky)

Lots of jobs could take advantage of this substrate

- $\triangleright$  Big compilations (compiling Chromium takes 16 hours on one core)
- $\triangleright$  Software test suites (unit tests, integration tests)
- **P** Ray-tracing (rendering one frame of a movie can take  $>12$  hours)
- $\blacktriangleright$  Video editing
- $\blacktriangleright$  Parallel jobs on large videos

### The gg intermediate representation

- $\blacktriangleright$  Types: values and thunks
- $\blacktriangleright$  Components
	- $\triangleright$  raw inputs ("V" value name or "T" thunk name)
	- $\triangleright$  forced inputs ("T" thunk name)
	- $\triangleright$  outputs (named byte vector, may be another thunk)
	- $\triangleright$  execution spec (e.g., Unix command line)
- $\blacktriangleright$  Addressing scheme
	- $\blacktriangleright$  "V" + hash of a byte vector
	- **If or** "T" + hash of a thunk's canonical representation + " $\#$ " + name of an output

### $\blacktriangleright$  Can express

- $\blacktriangleright$  Recursive fibonacci
- $\blacktriangleright$  Y combinator
- $\blacktriangleright$  Various everyday jobs

▶ Alpernas et al. (OOPSLA 2018): "Enforcing IFC policies is easy"

### Compilation



### Demo

### Compiling inkscape (600 kLOC)



### Compiling Chromium (24,000 kLOC)



### Tiny functions for lots of things. . .

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- Lepton: JPEG recompression
- ExCamera: video encoding with thousands of tiny tasks
- Salsify: real-time video with "functional" codec and transport
- gg: IR for "laptop to lambda" jobs with 8,000-way parallelism